

# Backdoor, Endianness, x86-64

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## CS 2130: Computer Systems and Organization 1

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## Announcements

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- Homework 4 due Monday after break on Gradescope
  - You have written most of this code already
  - Lab 7 may provide a fast way to get started
- Regrade requests for midterm 1 due Friday after break

## Backdoors

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**Backdoor:** secret way in to do new unexpected things

- Get around the normal barriers of behavior
- Ex: a way in to allow me to take complete control of your computer

**Exploit** - a way to use a vulnerability or backdoor that has been created

- Our exploit today: a **malicious payload**
  - A passcode and program
  - If it ever gets in memory, run my program regardless of what you want to do

## Ethics, Business, Tech

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Are there reasons to do this? Not to do this?

- No technical reason not to, it's easy to do!

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- Business implications (lawsuits, PR, etc) *Public relations.*

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Can we make a system where one bad actor can't break it?

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Can we make a system where one bad actor can't break it?

- Code reviews, double checks, verification systems, automated verification systems, ...

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Why does this work?

## Why?

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Why does this work?

- **It's all bytes!**
- Everything we store in computers are bytes
- We store code and data in the same place: memory

*Von Neumann model*

## It's all bytes

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Memory, Code, Data... It's all bytes!

- **Enumerate** - pick the meaning for each possible byte
- **Adjacency** - store bigger values together (sequentially)
- **Pointers** - a value treated as address of thing we are interested in

*You've seen all 3 of these already.*

## Enumerate

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**Enumerate** - pick the meaning for each possible byte

*Assign meaning to this byte.*

**What is 8-bit 0x54?**

Unsigned integer

Signed integer

Floating point w/ 4-bit exponent

ASCII

Bitvector sets

Our example ISA

eighty-four

positive eighty-four

twelve

capital letter T: T

The set {2, 3, 5}

Flip all bits of value in r1

## Adjacency

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**Adjacency** - store bigger values together (sequentially)

- An array: build bigger values out of many copies of the same type of small values

- Store them next to each other in memory

- Arithmetic to find any given value based on index

We know: ①. The address of the first element. ( $addr$ )

②. The index of that element ( $i$ )

Then the address of that element:  $addr + (i * \text{size\_of\_each\_element})$ .

## Adjacency

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**Adjacency** - store bigger values together (sequentially)

*One row in a database table, like one line in a CSV file.*

- Records, structures, classes

- Classes have fields! Store them adjacently

- Know how to access (add offsets from base address)

*addr + offset\_of -  $\pi$*

- If you tell me where object is, I can find fields

## Pointers

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**Pointers** - a value treated as address of thing we are interested in

- A value that really points to another value
- Easy to describe, hard to use properly
- We'll be talking about these a lot in this class!

## Pointers

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**Pointers** - a value treated as address of thing we are interested in

- Give us strange new powers (represent more complicated things), e.g.,
  - Variable-sized lists
  - Values that we don't know their type without looking
  - Dictionaries, maps

*Those 3 things, we combine them all together. And this is kind of how we're storing the data in the memory.*

## Programs Use These!

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How do our programs use these?

- Enumerated icodes, numbers
- Adjacenty stored instructions (PC+1)
- Pointers of where to jump/goto (addresses in memory)

## ToyISA Instructions

So far, only dealing with 8-bit machine!

→ 8-bit instructions/values/memory addresses

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		write $rA$ to memory at address $rB$
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$ For icode 6, increase $pc$ by 2 at end of instruction
7		Compare $rA$ as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment $pc$ as normal

## 64-bit Machines *The machine you have, is 64 bits.*

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64-bit machine: The registers are 64-bits

- i.e., r0, but also PC

Important to have large values. Why?

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Important to have large values. Why?

- Most important: PC and memory addresses *PC is wider* *We need space to do things.*

- How much memory could our 8-bit machine access?

*$2^8 = 256$  bytes. 00 - FF. Not a lot!*

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- Late 70s - 16 bits: **65536 Bytes** (*roughly 65 kilo bytes*).
- 80s - 32 bits:

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Important to have large values. Why?

- Most important: PC and memory addresses
- How much memory could our 8-bit machine access? **256 Bytes**
- Late 70s - 16 bits: **65536 Bytes**
- 80s - 32 bits:  $\approx 4$  billion bytes
- Today's processors - 64 bits:  $2^{64}$  addresses ( $2^{64}$  indices to memory)

## Aside: Powers of Two

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Value	base-10	Short form	Pronounced
$2^{10}$	$(10^3)1024$	Ki	Kilo
$2^{20}$	1,048,576	Mi	Mega
$2^{30}$	1,073,741,824	Gi	Giga
$2^{40}$	$(10^{12})1,099,511,627,776$	Ti	Tera
$2^{50}$	1,125,899,906,842,624	Pi	Peta
$2^{60}$	1,152,921,504,606,846,976	Ei	Exa

Example:  $2^{27}$  bytes

If I have 2 to some power, it works out to be roughly equivalent to a power of 10.

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Example:  $2^{27}$  bytes =  $2^7 \times 2^{20}$  bytes

$$2^m \times 2^n = 2^{m+n}$$

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Example:  $2^{27}$  bytes =  $2^7 \times 2^{20}$  bytes =  $2^7$  MiB = 128 MiB

## 64-bit Machines

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How much can we address with 64-bits?

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- 16 EiB ( $2^{64}$  addresses =  $2^4 \times \underline{2^{60}}$ )  
Exa

## 64-bit Machines

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How much can we address with 64-bits?

- 16 EiB ( $2^{64}$  addresses =  $2^4 \times 2^{60}$ )
- But I only have 8 GiB of RAM

I have 16 EiB of addresses. We can address more space we actually have.

But it could be used for virtual memory.

We will talk about it in CS0-2.

## A Challenge

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There is a disconnect:

- Registers: 64-bits values
- Memory: 8-bit values (i.e., **1 byte** values) *what we are storing is still 8 bits (1 byte),*
  - Each address addresses an 8-bit value in memory
  - Each address points to a 1-byte slot in memory

## A Challenge

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There is a disconnect:

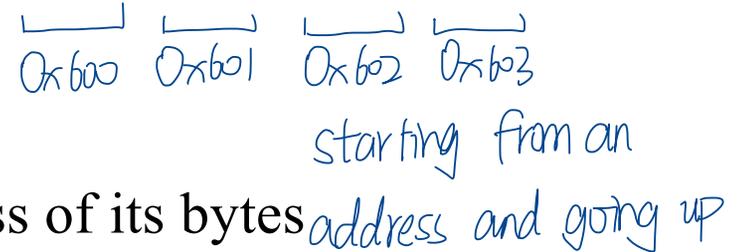
- Registers: 64-bits values
- Memory: 8-bit values (i.e., **1 byte** values)
  - Each address addresses an 8-bit value in memory
  - Each address points to a 1-byte slot in memory
- How do we store a 64-bit value in an 8-bit spot?

## Rules

0x|00|AB|CD|EF (4 bytes).

Rules to break “big values” into bytes (memory)

1. Break it into bytes
2. Store them adjacently
3. Address of the overall value = smallest address of its bytes
4. Order the bytes
  - If parts are ordered (i.e., array), first goes in smallest address
  - Else, hardware implementation gets to pick (!!)
  - Little-endian
  - Big-endian



## Ordering Values

0x|00|A B|C D|E F

### Little-endian

- Store the low order part/byte first
- Most hardware today is little-endian

$\underbrace{EF}_{0x00}$      $\underbrace{CD}_{0x01}$      $\underbrace{AB}_{0x02}$      $\underbrace{00}_{0x03}$

### Big-endian

- Store the high order part/byte first

$\underbrace{00}_{0x00}$      $\underbrace{AB}_{0x01}$      $\underbrace{CD}_{0x02}$      $\underbrace{EF}_{0x03}$

Why we want to talk about 2 ways?

Because people decided to do different things.

We write 00ABCDEF, but we calculate from F to 0,

Maybe that's the reason for processors to see EF first?

## Example

array of 2 numbers, each number should use 2 bytes.  
 Store [0x1234, 0x5678] at address 0xF00

	address	little endian	big endian
0x1234 {	0xF00	34	12
	0xF01	12	34
0x5678 {	0xF02	78	56
	0xF03	56	78

## Endianness

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Why do we study endianness?

- It is **everywhere**
- It is a source of weird bugs
- Ex: It's likely your computer uses:
  - Little-endian from CPU to memory
  - Big-endian from CPU to network
  - File formats are roughly half and half

*People didn't use the same thing.*

*In fact, your computers are probably doing different things now.*

Moving up!

## Assembly

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### General principle of all **assembly languages**

- Code (text, not binary!)
- 1 line of code = 1 machine instruction
- One-to-one reversible mapping between binary and assembly
  - We do not need to remember binary encodings!
  - A program will turn text to binary for us!

- ISA is like the grammar and vocabulary of a language.
- Assembly code is a sentence written in that language.

## Assembly

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### Features of assembly

- Automatic addresses - use **labels** to keep track of addresses
  - Assembler will remember location of labels and use where appropriate
  - Labels will not exist in machine code
- Metadata - data about data (*extra information*)
  - Data that helps turn assembly into code the machine can use
- As complicated as machine instructions
  - There are a lot of instructions, and it is one-to-one!

*It's going to replace them with the actual addresses when it builds the binary that we're going to run.*