

# Instruction Set Architectures, Stacks

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## CS 2130: Computer Systems and Organization 1

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## Announcements

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- Homework 3 due today
- Homework 4 available soon due Monday after Spring Break on Gradescope
  - You have written most of this code already

## Memory

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What kinds of things do we put in memory?

- Code: binary code like instructions in our example ISA

- Intel/AMD compatible: x86\_64
- Apple Mx and Ax, ARM: ARM
- And others!

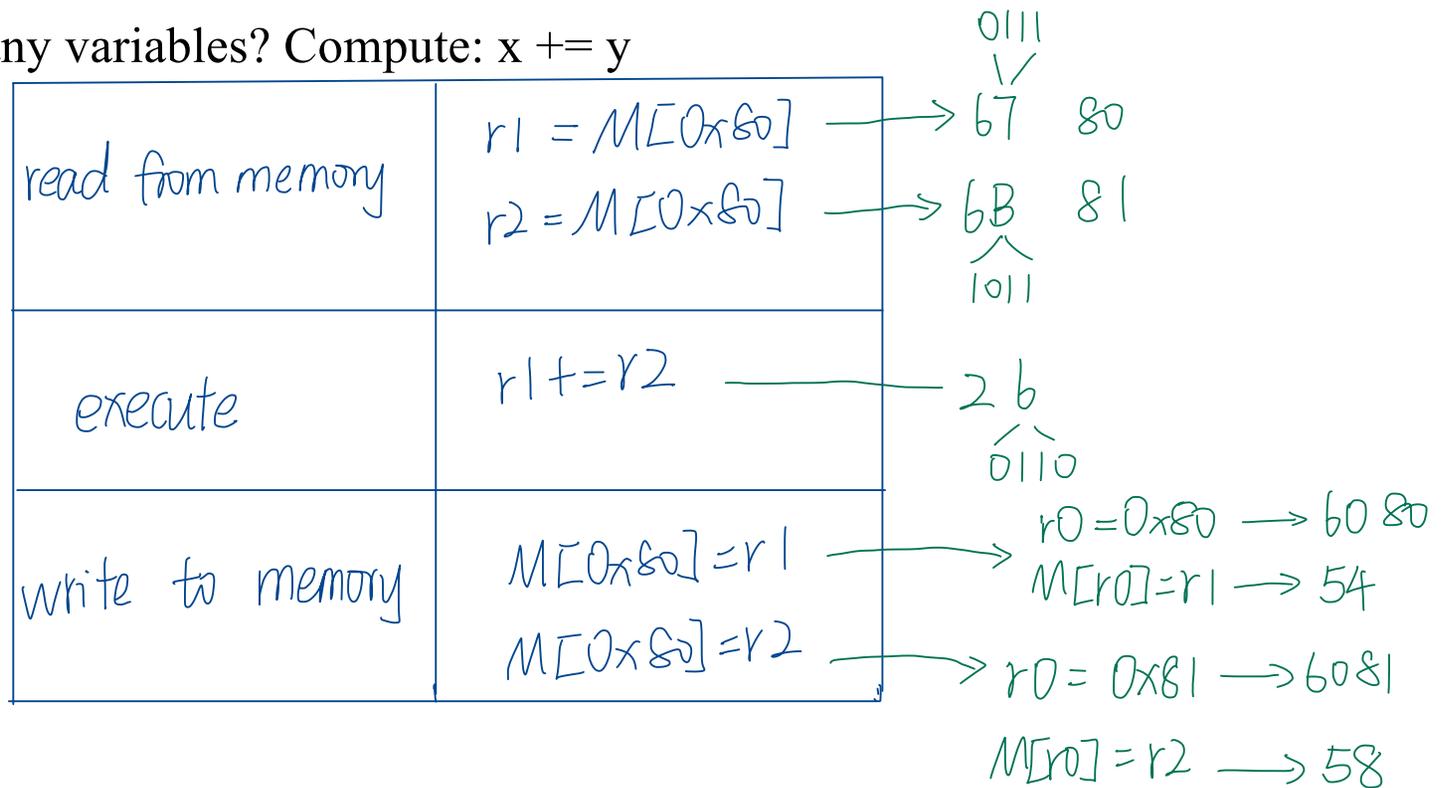
*Different CPUs have different ISAs.  
ISA describes how to understand  
the binary.*

- Variables: we may have more variables that will fit in registers
- Data Structures: organized data, collection of data
  - Arrays, lists, heaps, stacks, queues, ...

## Dealing with Variables and Memory

What if we have many variables? Compute:  $x += y$

$x = 0x80$   
 $y = 0x81$



67 80 6B 81 26 60 80 54 60 81 58

## Arrays

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**Array:** a sequence of values (collection of variables)

In Java, arrays have the following properties:

- Fixed number of values
- Not resizable
- All values are the same type

*they do have order  
index by integers.*

## Arrays

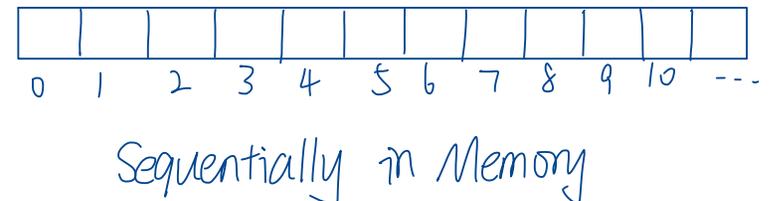
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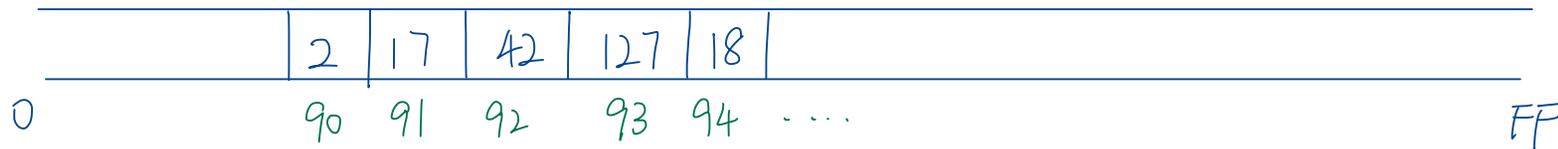
How do we store them in memory?



# Arrays

arr = {2, 17, 42, 127, 18 } @0x90

I assume all these are one byte so that each will fit in one of these slots.



$$\text{arr}[3] = \underbrace{0x90}_{\text{position of where our array start.}} + \underbrace{3}_{\text{index}} = 0x93$$

→ position of where our array start.

## Storing Arrays

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In memory, store array sequentially

- Pick address to store array
- Subsequent elements stored at following addresses
- Access elements with math

Example: Store array *arr* at **0x90**

- Access *arr*[3] as **0x90 + 3** assuming 1-byte values

$$arr[5] = 0x27$$

$$M[90+5] = 0x27$$

## What's Missing?

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What are we missing?

- Nothing says “this is an array” in memory
- Nothing says how long the array is

*Memory just stores a bunch of bytes.*

*They might be part of the array? Part of a large number? Part of an instruction? I don't know what they are. They're just values.*

*Some of our coding languages like Java, stores additional information about your arrays, like how long they are, what type they are.*

## Instructions

| icode | b | meaning  |
|-------|---|--|
| 0     |   | $rA = rB$  |
| 1     |   | $rA \&= rB$  |
| 2     |   | $rA += rB$   |
| 3     | 0 | $rA = \sim rA$   |
|       | 1 | $rA = !rA$   |
|       | 2 | $rA = -rA$   |
|       | 3 | $rA = pc$  |
| 4     |   | $rA =$ read from memory at address $rB$  |
| 5     |   | write $rA$ to memory at address $rB$   |
| 6     | 0 | $rA =$ read from memory at $pc + 1$  |
|       | 1 | $rA \&=$ read from memory at $pc + 1$  |
|       | 2 | $rA +=$ read from memory at $pc + 1$   |
|       | 3 | $rA =$ read from memory at the address stored at $pc + 1$  |
|       |   | For icode 6, increase $pc$ by 2 at end of instruction  |
| 7     |   | Compare $rA$ as 8-bit 2's-complement to 0<br>if $rA \leq 0$ set $pc = rB$<br>else increment $pc$ as normal |

## Instructions Set Architecture *We've designed an Instruction Set Architecture*

**Instruction Set Architecture (ISA)** is an abstract model of a computer defining how the CPU is controlled by software

- Conceptually, set of instructions that are possible and how they should be encoded
  - Results in many different machines to implement same ISA *Intel and AMD's CPU*
    - Example: How many machines implement our example ISA? *both using x86-64 ISA.*
  - Common in how we design hardware
- So they can run the same program.*

## Instructions Set Architecture

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**Instruction Set Architecture (ISA)** is an abstract model of a computer defining how the CPU is controlled by software

- Provides an abstraction layer between:
  - Everything computer is really doing (hardware)
  - What programmer using the computer needs to know (software)
- Hardware and Software engineers have freedom of design, if conforming to ISA
- Can change the machine without breaking any programs

*Lots of flexibility and freedom to build things that would be faster, like hyperthreading. I don't worry about on the software side.*

*Just make sure the code can be compiled to ISA. I can run it on hardware.*

## Instructions Set Architecture

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**Instruction Set Architecture (ISA)** is an abstract model of a computer defining how the CPU is controlled by software

- Provides an abstraction layer between:
  - Everything computer is really doing (hardware)
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*CSO: covering many of the times we'll need to think across this barrier*

*We're in general at this point, going to start staying just above this barrier and to the software side.*

## Instructions Set Architecture

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### *It's easy to handy* Backwards compatibility

0  
reserved

- Include flexibility to add additional instructions later *I can add more things.*
- Original instructions will still work *My new machine still can run the old code.  
old instructions still there.*
- Same program can be run on PC from 10+ years ago and new PC today

### Most manufacturers choose an ISA and stick with it

- Notable Exception: Apple  
*One manufacturer that has enough  
following to be able to make these changes and  
still come out ahead.*

*My new machine is much faster  
There may be some issues, but that is another story*

①. IBM's powerPC (1994 - 2005) MacOS 10.  
(PowerBook, PowerMac → PowerPC too hot/not scaling)

②. ⇒ Intel x86 (2006 - 2020)

RISC → CISC → RISC

- ① All programs had to be recompiled.
- ②. Hardware completely changed.
- ③. Backward compatibility became an issue.

Apple didn't want old programs to break. So they created: Rosetta

Rosetta is a translation layer that dynamically converted PowerPC instruction into x86 instructions. It essentially simulating the old ISA. Eventually, Apple "sunset" it and dropped support.

③. ⇒ Apple Silicon (2020 - now) M1 chip (their own ARM chip)

This is another ISA migration.

They resurrected Rosetta (Rosetta 2). It translates x86 → ARM on the fly.

- ①. Changing ISA is a massive ecosystem shift.
- ②. Most companies avoid changing ISA.
- ③. ISA stability enables software longevity.
- ④. Backward compatibility matters!
- ⑤. Hardware/software boundaries matter!
- ⑥. Binary translation preserves ecosystems.
- ⑦. ISA decisions affect decades of software.

## What about real ISAs?

CISC: Lots of instructions, some very specialized.

eg. x86-64: thousands of instructions, very complex.

RISC: fewer instructions, but each executes very efficiently.

eg. : RISC-V: 47 basic instructions. small and clean.

extensions: grow to hundreds of instructions.

ARM: 500 ~ 1000 instructions.

## Our Instructions Set Architecture

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What about our ISA?

- Enough instructions to compute what we need
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What about our ISA?

- Enough instructions to compute what we need
- As is, lot of things that are painful to do
  - This was on purpose! So we can see limitations of ISAs early
- Add any number of new instructions using the reserved bit (7)
- Missing something important: *Help to put variables in memory*

## Storing Variables in Memory

So far... we/compiler chose location for variable

Consider the following example:

$f(x)$ :

$a = x$

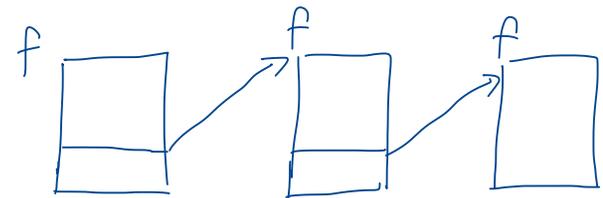
if  $(x \leq 0)$  return 0

else return  $f(x-1) + a$

$x = 5$

$a = 5$

$f(4) + 5 \Rightarrow$  Sums up the numbers between one and  $x$



### Recursion

- The formal study of a function that calls itself

while it computing the solution or the result for  $F$ , it actually calls itself in a smaller input.

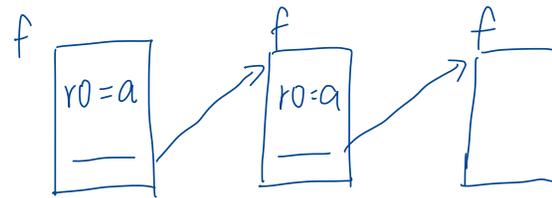
## Storing Variables in Memory

f(x):

a=x

if (x <= 0) return 0

else return f(x-1) + a



save it in a register R0? No!

save it to memory? M[0x80] No!

override this value over and over

When I call the same code again, jump to the same code again, I accidentally overwrite the value I really want.

Where do we store a?

We need something that will help us to organize memory so that we keep track of these variables.