

# Function Calls, Memory Instruction Set Architectures

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## CS 2130: Computer Systems and Organization 1

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## Announcements

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- Homework 3 due Next Monday on Gradescope
- Midterm 1 Friday in class
  - Written, closed notes
  - If you have SDAC, please schedule ASAP

## Exercises

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### Q6 Coding ToyISA

1 Point

In our example instruction set (Toy ISA) from class, encoding an operation may mean writing one or more instructions that collectively have exactly the same result as the operation we want. For example, to encode the operation `x = ~y`, we may encode it as `x = y; x = ~x;` (icode 0 then icode 3.0). However, while `y = ~y; x = y;` (icode 3.0 then icode 0) has the same effect, it would also modify y as well as x.

The source code operation `x = y + z;` could be implemented (efficiently using our instruction set as:

- One instruction
- Two instructions
- Three or more instructions
- It cannot be implemented with the Toy ISA instructions

## Exercises

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### Q7 ToyISA Encoding

2 Points

In our example instruction set (Toy ISA) from class, which of the following programs will compute `r0 = r0 - r2`?

- 3a 02
- 06 36 21
- 3a 12 00
- 06 34 21
- None of the above

Suppose we extended the ISA simulator you wrote in Lab 4 with the following code:

```
if (reserved == 1 && icode == 7) {  
    R[a] = R[b] >> M[oldPC + 1];  
    return oldPC + ____;  
}
```

1. [4 points] What is the value that we will need to increment the `oldPC`? (Fill in the blank below.)

```
return oldPC + _____;
```

2. [10 points] Using the new instruction above **at least once**, write a program that determines if both of the hexadecimal digits representing the byte in register 1 are odd, storing a 1 in register 0 if both are odd and a 0 if either are even. For example, if `r1` is `0x73`, store a `0x01` in `r0`; if `r1` is `0xE9`, store a `0x00` in `r0`. Do not change the values stored in `r1`, `r2`, or `r3`. Answer in hexadecimal bytes, separated by spaces. *Hint: You may need to write additional instructions.*

Answer: \_\_\_\_\_

3. [12 points] Complete the table below listing all the register values as hex digits after the following code executes. Assume that all registers start with value  $0\times 00$  and that the first instruction is at address  $0\times 00$ .

64 05 09 26 4E 62 F8 14 80

Register	Value
0	
1	
2	
3	

5. [15 points] Answer the following questions assuming 8-bit two's-complement numbers.

A. Compute the following sum, showing your work (such as carry bits, etc).

$$\begin{array}{r} 0\ 0\ 1\ 1\ 1\ 0\ 1\ 1 \\ +\ 0\ 1\ 1\ 0\ 0\ 1\ 1\ 1 \\ \hline \end{array}$$

B. Is your result a positive or negative number? (circle one)

**Positive**

**Negative**

C. Convert your solution to decimal.

Answer

7. [8 points] The following number is encoded as an 8-bit floating point number assuming a 4-bit exponent value.

01001111

A. Write the value in binary scientific notation.

Answer

B. Write the value in decimal.

Answer

10. [14 points] Using only 2-input `and`, `or`, `xor` gates, 1-input `not` gates, and constants 0 and/or 1, draw a circuit that has 3-bit input  $a$  with bits labeled  $a_2$ ,  $a_1$ , and  $a_0$ ; 3-bit output with bits labeled  $b_2$ ,  $b_1$ ,  $b_0$ ; which computes  $b = -a$  in 3-bit two's complement. *Clearly label your inputs and output bits individually.*

## High-level Instructions

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In general, 3 kinds of instructions

- **moves** - move values around without doing “work”
- **math** - broadly doing “work”
- **jumps** - jump to a new place in the code

## Memory

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What kinds of things do we put in memory?

- Code: binary code like instructions in our example ISA
  - Intel/AMD compatible: x86\_64
  - Apple Mx and Ax, ARM: ARM
  - And others!
- Variables: we may have more variables that will fit in registers
- Data Structures: organized data, collection of data
  - Arrays, lists, heaps, stacks, queues, ...

## Dealing with Variables and Memory

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What if we have many variables? Compute:  $x += y$

## Arrays

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**Array:** a sequence of values (collection of variables)

In Java, arrays have the following properties:

- Fixed number of values
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How do we store them in memory?

# Arrays

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## Storing Arrays

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In memory, store array sequentially

- Pick address to store array
- Subsequent elements stored at following addresses
- Access elements with math

Example: Store array *arr* at **0x90**

- Access *arr*[3] as **0x90 + 3** assuming 1-byte values

## What's Missing?

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What are we missing?

- Nothing says “this is an array” in memory
- Nothing says how long the array is

## Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		write $rA$ to memory at address $rB$
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$ For icode 6, increase $pc$ by 2 at end of instruction
7		Compare $rA$ as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment $pc$ as normal

## Instructions Set Architecture

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**Instruction Set Architecture (ISA)** is an abstract model of a computer defining how the CPU is controlled by software

- Conceptually, set of instructions that are possible and how they should be encoded
- Results in many different machines to implement same ISA
  - Example: How many machines implement our example ISA?
- Common in how we design hardware

## Instructions Set Architecture

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**Instruction Set Architecture (ISA)** is an abstract model of a computer defining how the CPU is controlled by software

- Provides an abstraction layer between:
  - Everything computer is really doing (hardware)
  - What programmer using the computer needs to know (software)
- Hardware and Software engineers have freedom of design, if conforming to ISA
- Can change the machine without breaking any programs

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  - What programmer using the computer needs to know (software)

*CSO: covering many of the times we'll need to think across this barrier*

## Instructions Set Architecture

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### Backwards compatibility

- Include flexibility to add additional instructions later
- Original instructions will still work
- Same program can be run on PC from 10+ years ago and new PC today

Most manufacturers choose an ISA and stick with it

- Notable Exception: Apple

## Our Instructions Set Architecture

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What about our ISA?

- Enough instructions to compute what we need
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## Our Instructions Set Architecture

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What about our ISA?

- Enough instructions to compute what we need
- As is, lot of things that are painful to do
  - This was on purpose! So we can see limitations of ISAs early
- Add any number of new instructions using the reserved bit (7)
- Missing something important: *Help to put variables in memory*

## Storing Variables in Memory

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So far... we/compiler chose location for variable

Consider the following example:

$f(x)$ :

$a=x$

if  $(x \leq 0)$  return 0

else return  $f(x-1) + a$

Recursion

- The formal study of a function that calls itself

## Storing Variables in Memory

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$f(x)$ :

$a=x$

if  $(x \leq 0)$  return 0

else return  $f(x-1) + a$

Where do we store  $a$ ?

## The Stack

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**Stack** - a last-in-first-out (LIFO) data structure

- The solution for solving this problem

**rsp** - Special register - the *stack* pointer

- Points to a special location in memory
- Two operations most ISAs support:
  - push - put a new value on the stack
  - pop - return the top value off the stack

## The Stack: Push and Pop

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push r0

- Put a value onto the “top” of the stack
  - $rsp -= 1$
  - $M[rsp] = r0$

pop r2

- Read value from “top”, save to register
  - $r2 = M[rsp]$
  - $rsp += 1$

## The Stack: Push and Pop

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## The Stack: Push and Pop

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## **What about real ISAs?**