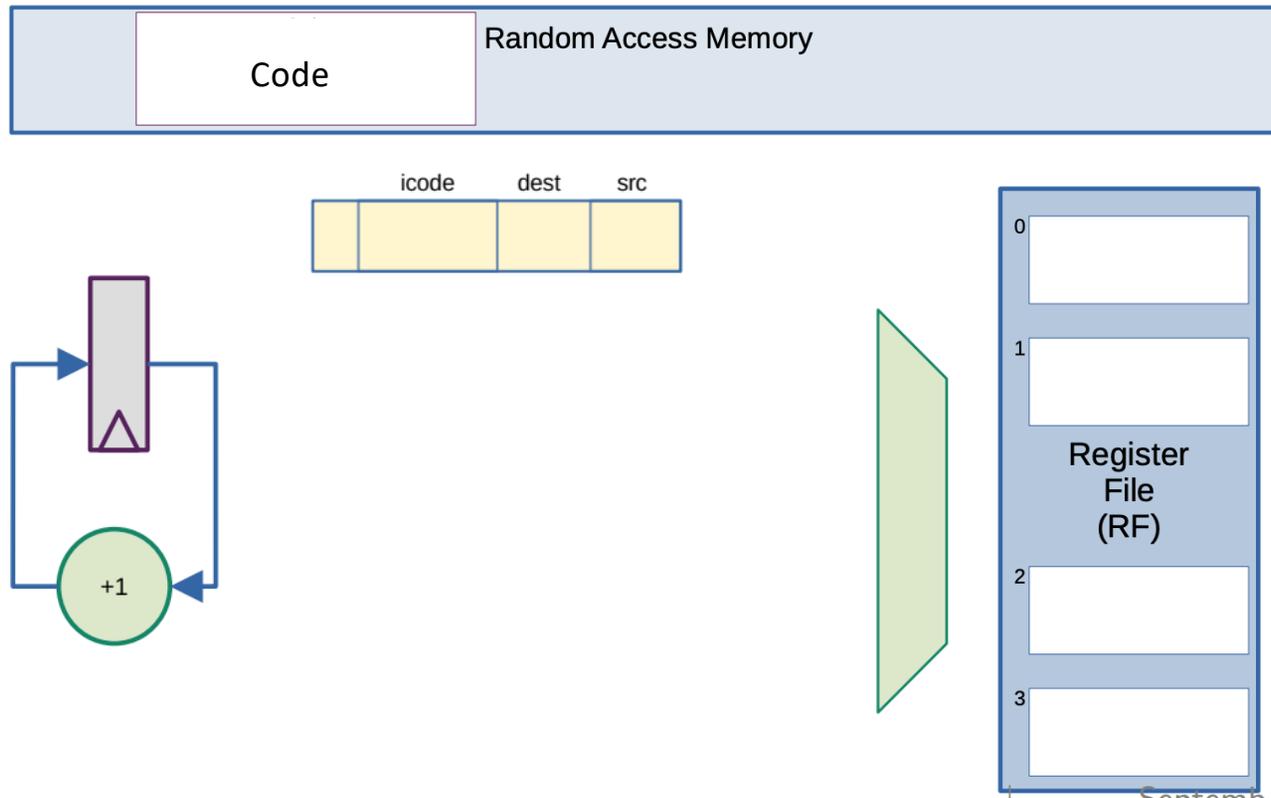


Fetch, Decode, Execute

CS 2130: Computer Systems and Organization 1

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Building a Computer



Question

What happens if we get the 0-byte instruction? 00

Our Computer's Instructions

Toy ISA 3-bit icode

icode	meaning
0	$rA = rB$
1	$rA \&= rB$
2	$rA += rB$
...	...
4	$rA =$ read from memory at address rB
5	write rA to memory at address rB
...	...
7	Compare rA as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment pc as normal

Our Computer's Instructions

Toy ISA 3-bit icode

icode	b	action
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$ For icode 6, increase pc by 2 at end of instruction

High-level Instructions

In general, 3 kinds of instructions

- **moves** - move values around without doing “work”
- **math** - broadly doing “work”
- **jumps** - jump to a new place in the code

Moves

Few forms

- Register to register (icode 0), $x = y$
- Register to/from memory (icodes 4-5), $x = M[b]$, $M[b] = x$

Memory

- Address: an index into memory.
 - Addresses are just (large) numbers
 - Usually we will not look at the number and trust it exists and is stored in a register

Moves

icode	b	action
0		$rA = rB$
3	3	$rA = pc$
4		$rA =$ read from memory at address rB
5		write rA to memory at address rB
6	0	$rA =$ read from memory at $pc + 1$
	3	$rA =$ read from memory at the address stored at $pc + 1$

Math

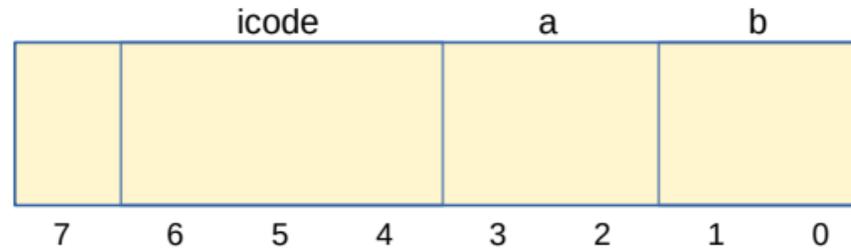
Broadly doing work

icode	b	meaning
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
6	1	$rA \&= \text{read from memory at pc} + 1$
	2	$rA += \text{read from memory at pc} + 1$

Note: We can implement other operations using these things!

icodes 3 and 6

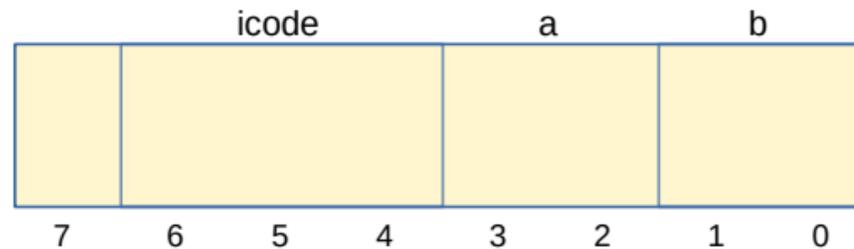
Special property of icodes 3 & 6: only one register used



icode	b	action
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$

icodes 3 and 6

Special property of icodes 3 & 6: only one register used



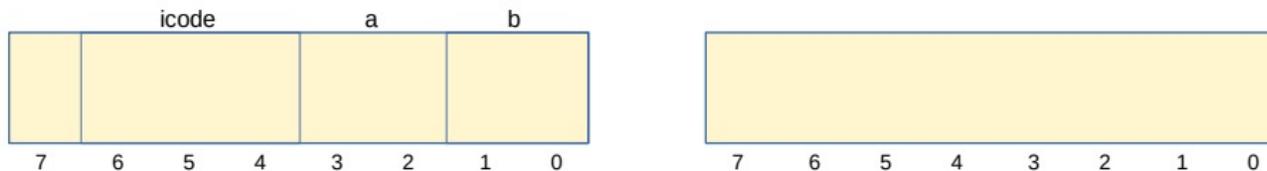
- Side effect: all bytes between 0 and 127 are valid instructions!
- As long as high-order bit is 0
- No syntax errors, any instruction given is valid

Immediate values

icode 6 provides literals, **immediate** values

icode	b	action
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$

For icode 6, increase pc by 2 at end of instruction



Encoding Instructions

Example 1: `r1 += 19`

Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA =$ read from memory at address rB
5		write rA to memory at address rB
6	0	$rA =$ read from memory at $pc + 1$
	1	$rA \&=$ read from memory at $pc + 1$
	2	$rA +=$ read from memory at $pc + 1$
	3	$rA =$ read from memory at the address stored at $pc + 1$ For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment pc as normal

Encoding Instructions

Example 2: $M[0x82] += r3$

Read memory at address 0x82, add r3, write back to memory at same address

Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		write rA to memory at address rB
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$
		For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment pc as normal

Jumps

- Moves and math are large portion of our code
- We also need **control constructs**
 - Change what we are going to do next
 - if, while, for, functions, ...
- Jumps provide mechanism to perform these control constructs
- We jump by assigning a new value to the program counter PC

Jumps

- For example, consider an if

Jumps

icode	meaning
7	Compare rA as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment pc as normal

Instruction icode 7 provides a **conditional** jump

- Real code will also provide an **unconditional** jump, but a conditional jump is sufficient

Writing Code

We can now write any* program!

- When you run code, it is being turned into instructions like ours
- Modern computers use a larger pool of instructions than we have (we will get there)

*we do have some limitations, since we can only represent 8-bit values and some operations may be tedious.

Our code to this machine code

How do we turn our control constructs into jump statements?

if/else to jump

while to jump

Function Calls

Encoding Instructions

Example 3: if r0 < 9 jump to 0x42

Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA =$ read from memory at address rB
5		write rA to memory at address rB
6	0	$rA =$ read from memory at $pc + 1$
	1	$rA \&=$ read from memory at $pc + 1$
	2	$rA +=$ read from memory at $pc + 1$
	3	$rA =$ read from memory at the address stored at $pc + 1$
		For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment pc as normal