



Toy Instruction Set Architecture

CS 2130: Computer Systems and Organization 1
February 13, 2026

Announcements

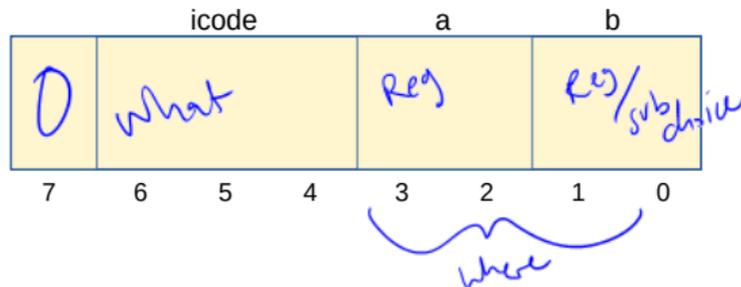
→ Feb 23

- Homework 3 due Monday on Gradescope
- Midterm 1 next Friday (February 20, 2026) in class
 - Written, closed notes
 - If you have SDAC, please schedule ASAP
- Review session in class next Wednesday

Encoding Instructions

Encoding of Instructions

- 3-bit icode (which operation to perform)
 - Numeric mapping from icode to operation
- Which registers to use (2 bits each)
- Reserved bit for future expansion



High-level Instructions

In general, 3 kinds of instructions

- **moves** - move values around without doing “work”
- **math** - broadly doing “work”
- **jumps** - jump to a new place in the code

Moves

icode	b	action
0		$rA = rB$
3	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		$\text{write } rA \text{ to memory at address } rB$
6	0	$rA = \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$

Math

Broadly doing work

icode	b	meaning
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
6	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$

Note: We can implement other operations using these things!

Encoding Instructions

Example 2: M[0x~~2~~^{1A}] += r3

Read memory at address 0x~~2~~^{1A}, add r3, write back to memory at same address

Instructions

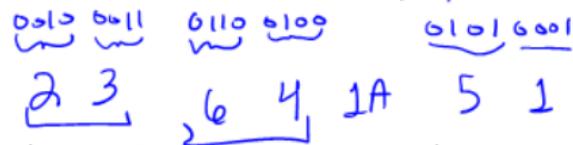
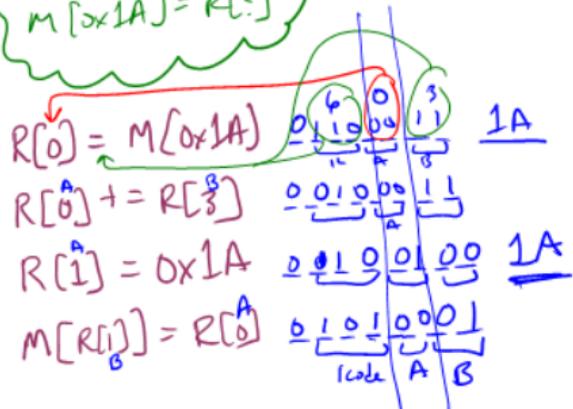
icode	b	meaning
0		rA = rB
1		rA &= rB
2		rA += rB
3	0	rA = ~rA
	1	rA = !rA
	2	rA = -rA
	3	rA = pc
→ 4		rA = read from memory at address rB
5		write rA to memory at address rB
→ 6	0	rA = read from memory at pc + 1
	1	rA &= read from memory at pc + 1
	2	rA += read from memory at pc + 1
→ 3		rA = read from memory at the address stored at pc + 1 For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0 if rA ≤ 0 set pc = rB else increment pc as normal

63 1A 23 64 1A 51

0110 0011
 6 3 1A

∴ M[0x1A] += R[3]

R[?] = M[0x1A]
 R[?] += R[3]
 M[0x1A] = R[?]

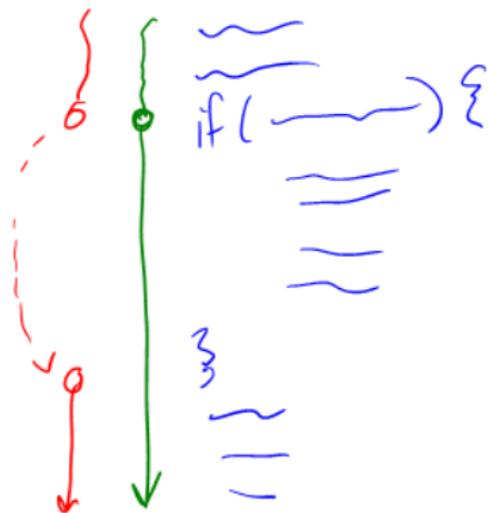


Jumps

- Moves and math are large portion of our code
- We also need **control constructs**
 - Change what we are going to do next
 - `if`, `while`, `for`, functions, ...
- Jumps provide mechanism to perform these control constructs
- We jump by assigning a new value to the program counter PC

Jumps

For example, consider an `if`



Jumps

icode	meaning
7	Compare rA as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment pc as normal

Instruction icode 7 provides a **conditional** jump

- Real code will also provide an **unconditional** jump, but a conditional jump is sufficient

Writing Code

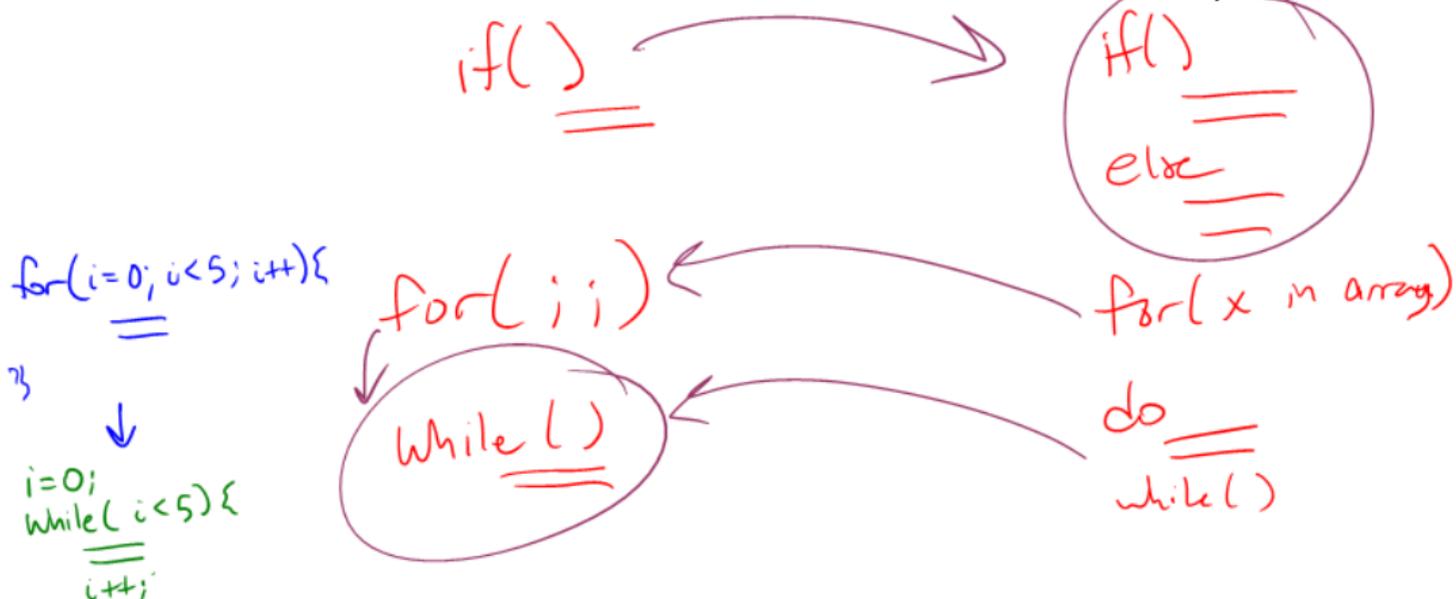
We can now write any* program!

- When you run code, it is being turned into instructions like ours
- Modern computers use a larger pool of instructions than we have (we will get there)

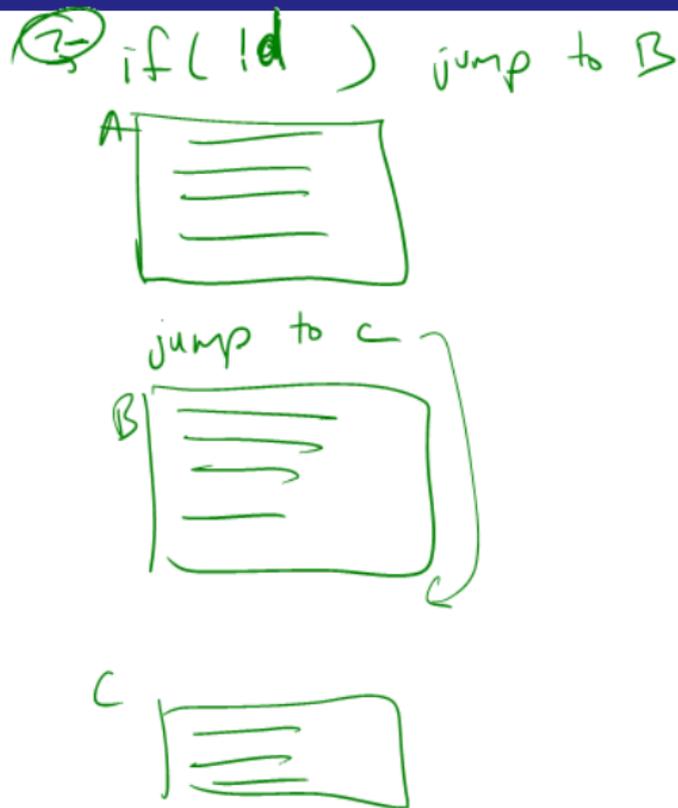
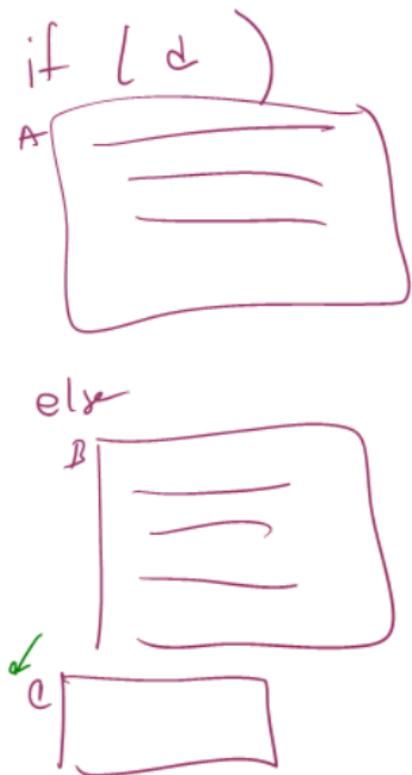
*we do have some limitations, since we can only represent 8-bit values and some operations may be tedious.

Our code to this machine code

How do we turn our control constructs into jump statements?

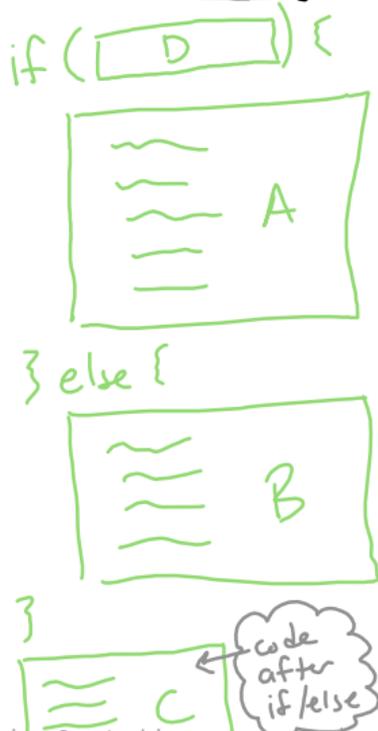


if/else to jump



if/else to jump

Pseudocode using if/else

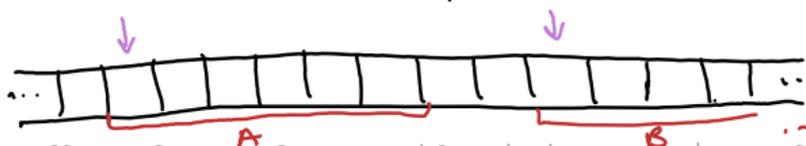


Notes!

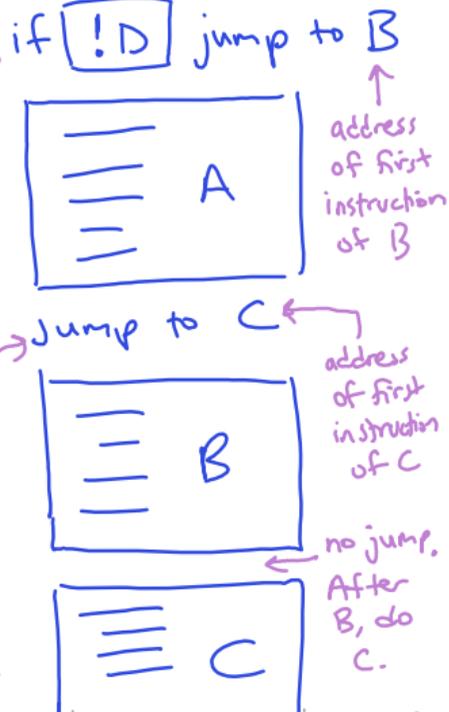
if D is true:
run code in A, then C
→ skip B

if D is false:
run code in B, then C
→ skip A

Code is in memory (think array)



Using Jumps



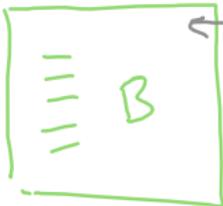
while to jump

Pseudocode of while loop:

while (C) {



}



Code after loop

Notes:

→ if C is true, run code in A, then go back and check C again

if C is false, skip A, go to B

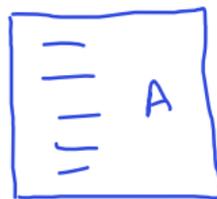
We have two options!

- jump to the check of C

- check at end of A before jumping back

Option 1

D → if (!C) jump to B



jump to D unconditionally

address of first instruction of ... ↓

Option 2

if (!C) jump to B



if (C) jump to A

