

Toy Instruction Set Architecture

aka Writing Programs for our computer

CS 2130: Computer Systems and Organization 1

February 11, 2026

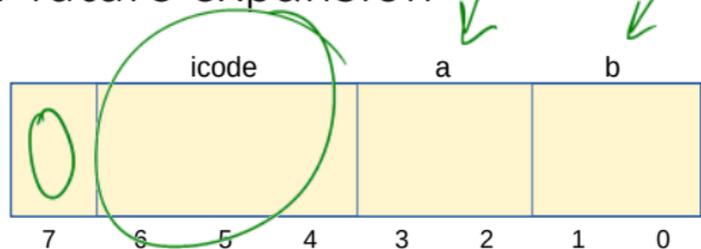
Announcements

- Homework 3 due Monday on Gradescope
- Midterm 1 next Friday (February 20) in class
 - Written, closed notes
 - If you have SDAC, please schedule ASAP

Encoding Instructions

Encoding of Instructions

- 3-bit icode (which operation to perform)
 - Numeric mapping from icode to operation
- Which registers to use (2 bits each) $R[a]$ $R[b]$
- Reserved bit for future expansion



Toy ISA Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA =$ read from memory at address rB
5		write rA to memory at address rB
6	0	$rA =$ read from memory at $pc + 1$
	1	$rA \&=$ read from memory at $pc + 1$
	2	$rA +=$ read from memory at $pc + 1$
	3	$rA =$ read from memory at the address stored at $pc + 1$
		For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment pc as normal

High-level Instructions

In general, 3 kinds of instructions

- **moves** - move values around without doing “work”
- **math** - broadly doing “work”
- **jumps** - jump to a new place in the code

Moves

Few forms

- Register to register (icode 0), $x = y$
- Register to/from memory (icodes 4-5), ~~$x = M[b]$, $M[a] = x$~~

$$R[a] = R[b]$$

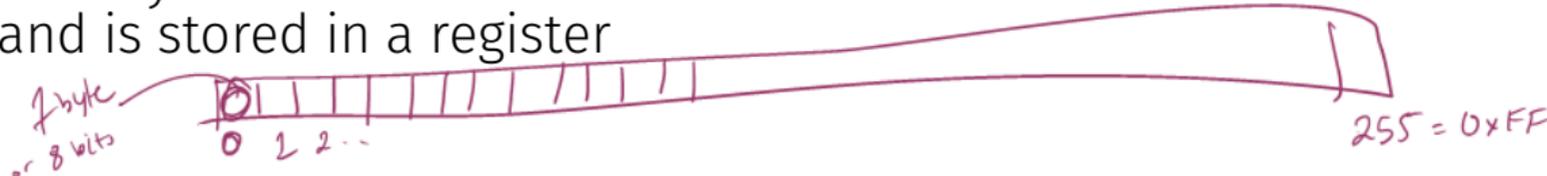
$$b = 2 \text{ bits}$$
$$R[b] = 8 \text{ bits}$$

Memory

- **Address:** an index into memory.
 - Addresses are just (large) numbers
 - Usually we will not look at the number and trust it exists and is stored in a register

$$R[a] = M[R[b]]$$

$$M[R[b]] = R[a]$$



Moves

icode	b	action
0		$rA = rB$
3	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		$\text{write } rA \text{ to memory at address } rB$
6	0	$rA = \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$

Math

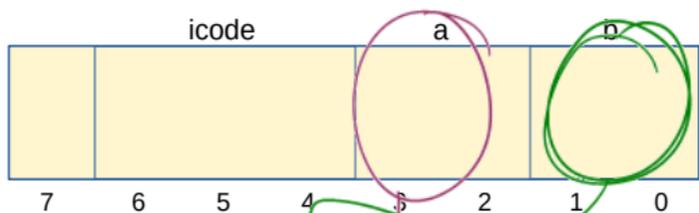
Broadly doing work

icode	b	meaning
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
6	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$

Note: We can implement other operations using these things!

icodes 3 and 6

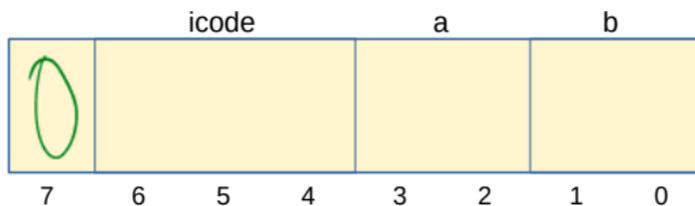
Special property of icodes 3 & 6: only one register used



icode	b	action
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$

icodes 3 and 6

Special property of 3 & 6: only one register used



- Side effect: all bytes between 0 and 127 are valid instructions!
- As long as high-order bit is 0
- No syntax errors, any instruction given is valid

Immediate values

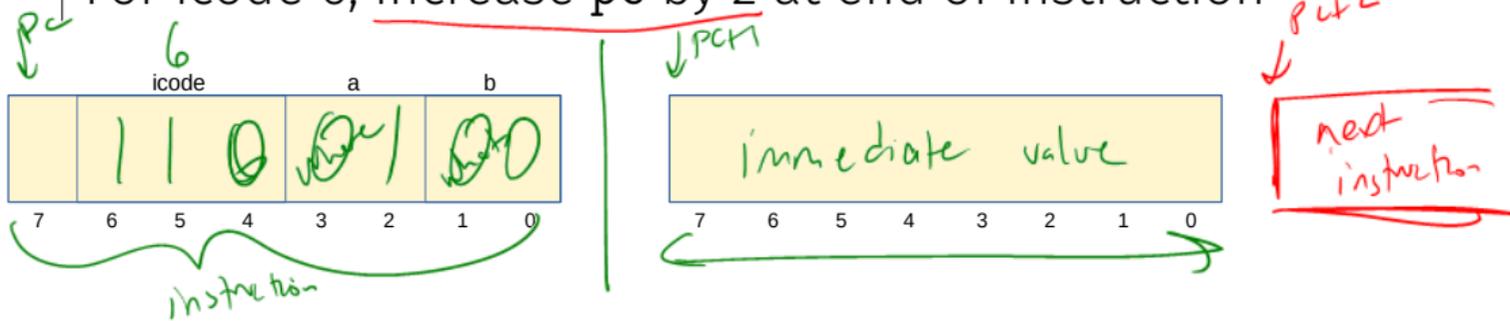
icode 6 provides literals, **immediate** values

$x = 25$
 $int\ y = 3;$

icode	b	action
6	0	rA = read from memory at pc + 1
	1	rA &= read from memory at pc + 1
	2	rA += read from memory at pc + 1
	3	rA = read from memory at the address stored at pc + 1

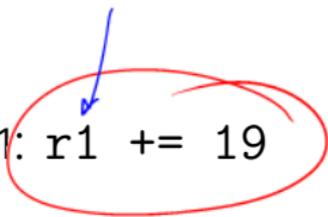
For icode 6, increase pc by 2 at end of instruction

$$r[A] = M[M[pc+1]]$$



Encoding Instructions

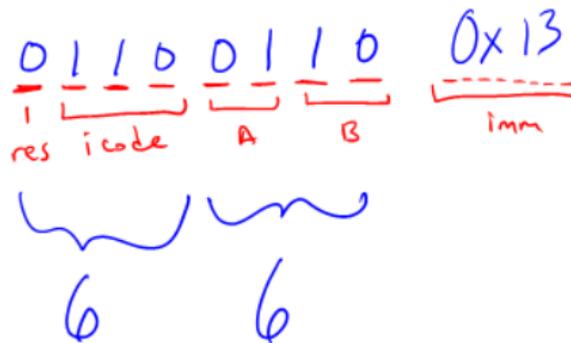
Example 1: `r1 += 19`



Instructions

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1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA =$ read from memory at address rB
5		write rA to memory at address rB
6	0	$rA =$ read from memory at $pc + 1$
	1	$rA \&=$ read from memory at $pc + 1$
	2	$rA +=$ read from memory at $pc + 1$
	3	$rA =$ read from memory at the address stored at $pc + 1$ For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment pc as normal

$R[1] += 19$



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Encoding Instructions

Example 2: `M[0x82] += r3`

Read memory at address `0x82`, add `r3`, write back to memory at same address

Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA =$ read from memory at address rB
5		write rA to memory at address rB
6	0	$rA =$ read from memory at $pc + 1$
	1	$rA \&=$ read from memory at $pc + 1$
	2	$rA +=$ read from memory at $pc + 1$
	3	$rA =$ read from memory at the address stored at $pc + 1$
		For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment pc as normal

Writing Code: Homework Hints

1. Write pseudocode that does the desired task
- 2-3 ... deal with control flow
4. Split multi-operation lines into series of single-operation lines
 $x = y - z$; becomes $x = y$; $x -= z$;
5. Convert operations to those in our instruction set
 $x -= z$; becomes $w = z$; $w = -w$; $x += w$;
6. ... deal with loops
7. Assign variables to our four registers, ex: $r0=x$, $r1=y$, $r2=z$, $r3=w$
 $r0 = r1$; $r3 = r2$; $r3 = -r3$; $r0 += r3$
- 10- Write those instructions into triples, then hex