

Function Calls, Memory Instruction Set Architectures

CS 2130: Computer Systems and Organization 1

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Announcements

- Homework 3 due Wednesday at 11:59pm on Gradescope
- Midterm 1 Friday (October 3, 2025) in class
 - Written, closed notes
 - If you have SDAC, please schedule ASAP

Quiz Questions – Quiz 1

Q5.3 XOR

1 Point

Suppose we then shift it back and xor it with the original, like

```
((0xCA >> 3) << 3) ^ 0xCA.
```

The result is:

- the same for both signed and unsigned integers
- larger for signed than unsigned integers
- larger for unsigned than signed integers
- O there is no way to know



Quiz Questions – Quiz 3

Q8 Counter

1 Point

To build a 4-bit counter circuit, we could directly connect the output of the increment circuit back to the input.

True

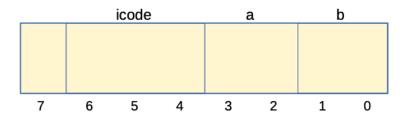
False



Encoding Instructions

Encoding of Instructions

- 3-bit icode (which operation to perform)
 - Numeric mapping from icode to operation
- Which registers to use (2 bits each)
- Reserved bit for future expansion





High-level Instructions

In general, 3 kinds of instructions

- moves move values around without doing "work"
- math broadly doing "work"
- **jumps** jump to a new place in the code



Memory

What kinds of things do we put in memory?

- Code: binary code like instructions in our example ISA
 - Intel/AMD compatible: x86 64
 - Apple Mx and Ax, ARM: ARM
 - And others!
- Variables: we may have more variables that will fit in registers
- Data Structures: organized data, collection of data
 - Arrays, lists, heaps, stacks, queues, ...



Dealing with Variables and Memory

What if we have many variables? Compute: x += y



Arrays

Array: a sequence of values (collection of variables)

In Java, arrays have the following properties:

- Fixed number of values
- Not resizable
- All values are the same type



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How do we store them in memory?



Arrays



Storing Arrays

In memory, store array sequentially

- Pick address to store array
- Subsequent elements stored at following addresses
- Access elements with math

Example: Store array *arr* at **0x90**

• Access arr[3] as 0x90 + 3 assuming 1-byte values



What's Missing?

What are we missing?

- Nothing says "this is an array" in memory
- Nothing says how long the array is

Instructions

icode	b	meaning
0		rA = rB
1		rA &= rB
2		rA += rB
3	0	rA = ~rA
	1	rA = !rA
	2	rA = -rA
	3	rA = pc
4		rA = read from memory at address rB
5		write rA to memory at address rB
6	0	rA = read from memory at pc + 1
	1	rA &= read from memory at pc + 1
	2	$\mathtt{rA} += \mathtt{read}$ from memory at $\mathtt{pc} + \mathtt{1}$
	3	rA = read from memory at the address stored at pc + 1
		For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0
		if $rA \le 0$ set $pc = rB$
		else increment pc as normal



Instruction Set Architecture (ISA) is an abstract model of a computer defining how the CPU is controlled by software

- Conceptually, set of instructions that are possible and how they should be encoded
- Results in many different machines to implement same ISA
 - Example: How many machines implement our example ISA?
- Common in how we design hardware



Instruction Set Architecture (ISA) is an abstract model of a computer defining how the CPU is controlled by software

- Provides an abstraction layer between:
 - Everything computer is really doing (hardware)
 - What programmer using the computer needs to know (software)
- Hardware and Software engineers have freedom of design, if conforming to ISA
- Can change the machine without breaking any programs



Instruction Set Architecture (ISA) is an abstract model of a computer defining how the CPU is controlled by software

- Provides an abstraction layer between:
 - Everything computer is really doing (hardware)
 - What programmer using the computer needs to know (software)

CSO: covering many of the times we'll need to think across this barrier



Backwards compatibility

- Include flexibility to add additional instructions later
- Original instructions will still work
- Same program can be run on PC from 10+ years ago and new PC today

Most manufacturers choose an ISA and stick with it

• Notable Exception: Apple



What about our ISA?

- Enough instructions to compute what we need
- As is, lot of things that are painful to do
 - This was on purpose! So we can see limitations of ISAs early



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What about our ISA?

- Enough instructions to compute what we need
- As is, lot of things that are painful to do
 - This was on purpose! So we can see limitations of ISAs early
- Add any number of new instructions using the reserved bit (7)
- Missing something important: *Help to put variables in memory*

Storing Variables in Memory

So far... we/compiler chose location for variable

```
Consider the following example:

f(x):

a=x

if (x <= 0) return 0

else return f(x-1) + a
```

Recursion

• The formal study of a function that calls itself



Storing Variables in Memory

```
f(x):
    a=x
    if (x <= 0) return 0
    else return f(x-1) + a</pre>
```

Where do we store a?



The Stack

Stack - a last-in-first-out (LIFO) data structure

• The solution for solving this problem

rsp - Special register - the *stack* pointer

- Points to a special location in memory
- Two operations most ISAs support:
 - push put a new value on the stack
 - pop return the top value off the stack

The Stack: Push and Pop

push r0

- Put a value onto the "top" of the stack
 - rsp -= 1
 - M[rsp] = r0

pop r2

- Read value from "top", save to register
 - r2 = M[rsp]
 - rsp += 1



The Stack: Push and Pop



The Stack: Push and Pop



What about real ISAs?