

Toy Instruction Set Architecture

CS 2130: Computer Systems and Organization 1

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Announcements

- Homework 3 due Monday at 11:59pm on Gradescope
- Quiz 4 available today, due Sunday at 11:59pm
- Midterm 1 next Friday (October 3, 2025) in class
 - Written, closed notes
 - If you have SDAC, please schedule ASAP

High-level Instructions

In general, 3 kinds of instructions

- **moves** - move values around without doing “work”
- **math** - broadly doing “work”
- **jumps** - jump to a new place in the code

Encoding Instructions

idea: ①. I have a value in memory at address hex 82

Example 2: $M[0x82] += r3$

②. I want to add whatever in R3 to that value.

Read memory at address 0x82, add r3, write back to memory at same address

One point: No instructions allow us to pass an immediate value as the address.

So let's save ourselves some time: just put 82 in a register.

Then we use icode 4 to read it out, icode 5 to write it back.

$r2 = 0x82$	$\begin{array}{c} 0 \ 110 \ 10 \ 00 \\ \hline 0 \ 6 \ 100 \ 01 \ 10 \\ \hline 0 \ 4 \ 010 \ 01 \ 11 \\ \hline 0 \ 2 \ 101 \ 01 \ 10 \\ \hline 0 \ 5 \end{array}$
$r1 = M[r2]$	
$r1 += r3$	
$M[r2] = r1$	

82

our machine reads 0 and 1.

But, for us \rightarrow easier to read \rightarrow pairs of hex

68 82 46 27 5b

Encoding Instructions

icode	b	meaning
0		rA = rB
1		rA &= rB
2		rA += rB
3	0	rA = ~rA
	1	rA = !rA
	2	rA = -rA
	3	rA = pc
4		rA = read from memory at address rB
5		write rA to memory at address rB
6	0	rA = read from memory at pc + 1
	1	rA &= read from memory at pc + 1
	2	rA += read from memory at pc + 1
	3	rA = read from memory at the address stored at pc + 1
7		For icode 6, increase pc by 2 at end of instruction
		Compare rA as 8-bit 2's-complement to 0
		if rA <= 0 set pc = rB
		else increment pc as normal

Example 3: if r0 < 9 jump to 0x42

I don't have an instruction say $r0 < 9$.
I need " $r0 \leq 0$ " for icode 7, what should I do?

$$r0 < 9 \Leftrightarrow r0 \leq 8 \Leftrightarrow (r0 - 8) \leq 0$$

$$\Leftrightarrow r0 += -8 \text{ (0xF8)}$$

$$r0 \leq 0$$

$$r1 = 0x42 \quad \begin{array}{r} 0 \ 110 \ 01 \ 00 \\ 6 \quad 4 \quad 42 \end{array}$$

$$r0 += F8 \quad \begin{array}{r} 0 \ 110 \ 00 \ 10 \\ 6 \quad 2 \quad F8 \end{array}$$

$$\text{if } r0 \leq 0, PC = r1 \quad \begin{array}{r} 0 \ 111 \ 00 \ 01 \\ 7 \quad 1 \end{array}$$

6442 62F871

Function Calls

Dealing with Variables and Memory

What if we have many variables? Compute: $x += y$

Arrays

Array: a sequence of values (collection of variables)

In Java, arrays have the following properties:

- Fixed number of values
- Not resizable
- All values are the same type

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How do we store them in memory?

Arrays

Storing Arrays

In memory, store array sequentially

- Pick address to store array
- Subsequent elements stored at following addresses
- Access elements with math

Example: Store array *arr* at **0x90**

- Access *arr*[3] as **0x90 + 3** assuming 1-byte values

Quiz Questions – Quiz 1

Q5.3 XOR

1 Point

Suppose we then shift it back and xor it with the original, like

```
((0xCA >> 3) << 3) ^ 0xCA.
```

The result is:

- ☒ the same for both signed and unsigned integers
- ☐ larger for signed than unsigned integers
- ☐ larger for unsigned than signed integers
- ☐ there is no way to know

Quiz Questions – Quiz 2

Q4 Floating Point

2 Points

Assume we will use 8-bit floating-point numbers with **3 fraction bits**. How would we encode the binary number `-010110000` into this 8-bit floating point representation?

- ☐ `0 011 1110`
- ☐ `1 011 1110`
- ☐ `0 0111 011`
- ☐ `0 1110 011`
- ☐ `1 0111 011`
- ☒ `1 1110 011`

Quiz Questions – Quiz 3

Q3 Coding hardware

1 Point

When coding in a hardware description language (code that can be turned into circuits), there are no typical control constructs like `if`, `while`, and `for`; in addition, which of the following is **not** permitted?

- ☐ Accessing the same variable twice, like `y = x + 1; z = x - 1;`
- ☒ Assigning to the same variable twice, like `y = x + 1; y = w - z;`
- ☐ Conditional operations, like `y = (x < 0) ? x : -x;`
- ☐ Including several operators in a single expression, like `y = (x + y) ^ z;`

Quiz Questions – Quiz 3

Q5 Cycles

1 Point

In class, we built a computer that we could program with 1-byte instructions, containing an icode, source, and destination registers. In each cycle, the logic circuits **only** calculated the operation specified by the given icode, which would be later written to a register.

☐ True

☒ False

Quiz Questions – Quiz 3

Q6 Coding ToyISA

1 Point

In our example instruction set (Toy ISA) from class, encoding an operation may mean writing one or more instructions that collectively have exactly the same result as the operation we want. For example, to encode the operation `x = ~y`, we may encode it as `x = y; x = ~x;` (icode 0 then icode 3.0). However, while `y = ~y; x = y;` (icode 3.0 then icode 0) has the same effect, it would also modify y as well as x.

The source code operation `x = y + z;` could be implemented (efficiently using our instruction set as:

- ☐ One instruction
- ☒ Two instructions
- ☐ Three or more instructions
- ☐ It cannot be implemented with the Toy ISA instructions

Quiz Questions – Quiz 3

Q7 ToyISA Encoding

2 Points

In our example instruction set (Toy ISA) from class, which of the following programs will compute `r0 = r0 - r2`?

- ☐ 3a 02
- ☒ 06 36 21
- ☐ 3a 12 00
- ☐ 06 34 21
- ☐ None of the above

Quiz Questions – Quiz 3

Q8 Counter

1 Point

To build a 4-bit counter circuit, we could directly connect the output of the increment circuit back to the input.

- ☐ True
- ☒ False