

# Building to a Computer

## Fetch, Decode, Execute

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### CS 2130: Computer Systems and Organization 1

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## Announcements

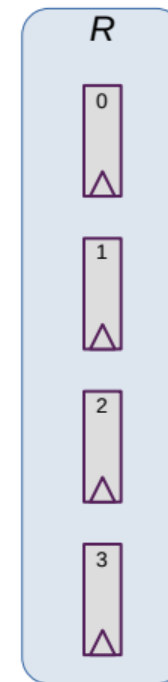
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- Quiz 3 available today, due Sunday by 11:59pm
- Homework 2 due Monday

## Writing

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$R[j] = y$  - connect  $y$  to input of registers based on index  $j$



## Code

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How do we run code? What do we need?

Consider the following code:

...

8:  $x = 16$

9:  $y = x$

10:  $x += y$

...

What is the value of  $x$  after line 10?

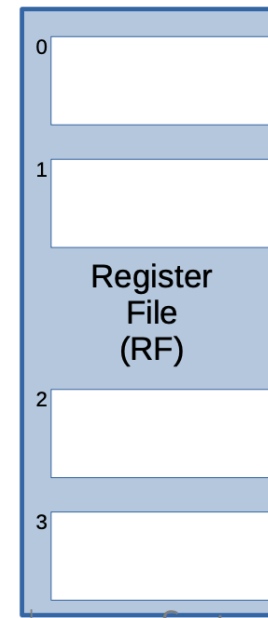
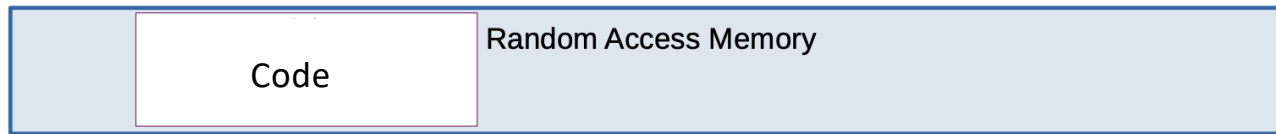
## Bookkeeping

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What do we need to keep track of?

- **Code** - the program we are running
  - RAM (Random Access Memory)
- **State** - things that may change value (i.e., variables)
  - Register file - can read and write values each cycle
- **Program Counter (PC)** - where we are in our code
  - Single register - byte number in memory for next instruction

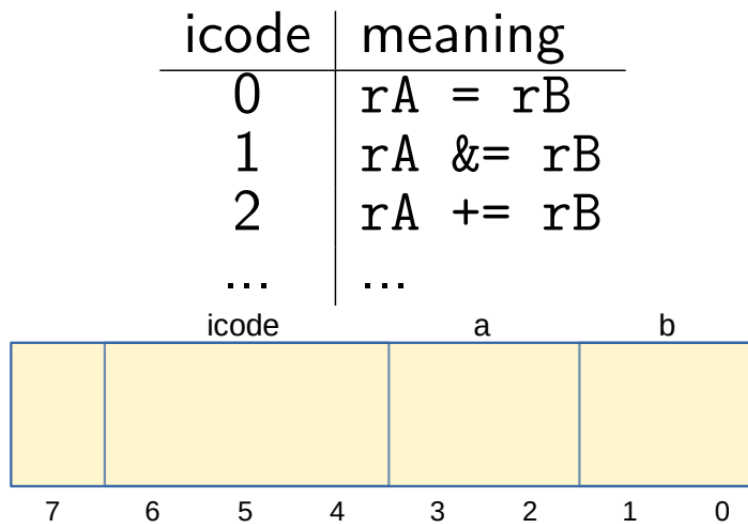
# Building a Computer



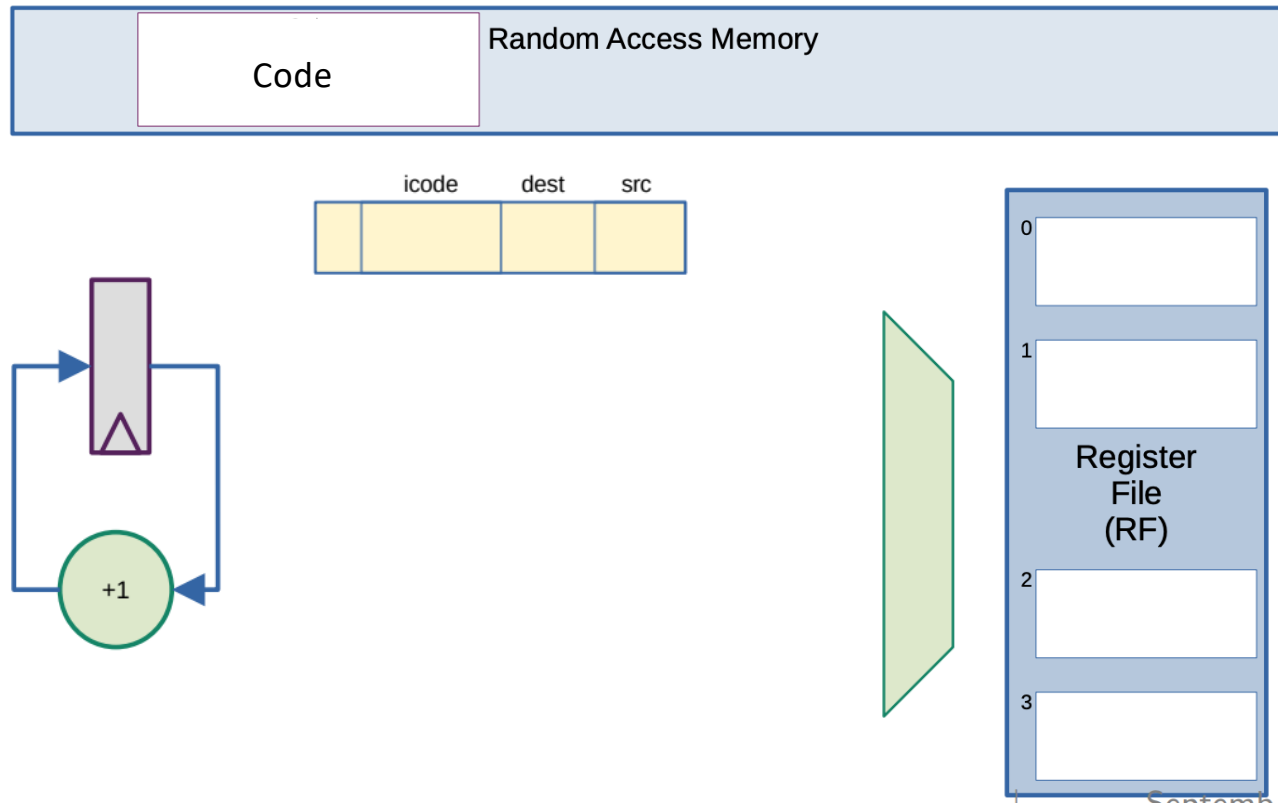
## Encoding Instructions

### Encoding of Instructions (**icode** or **opcode**)

- Numeric mapping from icode to operation



# Building a Computer





## Question

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What happens if we get the 0-byte instruction? 00

## Our Computer's Instructions

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Toy ISA 3-bit icode

icode	meaning
0	$rA = rB$
1	$rA \&= rB$
2	$rA += rB$
...	...
4	$rA = \text{read from memory at address } rB$
5	write $rA$ to memory at address $rB$
...	...
7	Compare $rA$ as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment $pc$ as normal

## Our Computer's Instructions

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Toy ISA 3-bit icode

icode	b	action
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$
For icode 6, increase pc by 2 at end of instruction		

## High-level Instructions

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In general, 3 kinds of instructions

- **moves** - move values around without doing “work”
- **math** - broadly doing “work”
- **jumps** - jump to a new place in the code

## Moves

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Few forms

- Register to register (icode 0),  $x = y$
- Register to/from memory (icodes 4-5),  $x = M[b]$ ,  $M[b] = x$

Memory

- Address: an index into memory.
  - Addresses are just (large) numbers
  - Usually we will not look at the number and trust it exists and is stored in a register

## Moves

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icode	b	action
0		$rA = rB$
3	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		$\text{write } rA \text{ to memory at address } rB$
6	0	$rA = \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$

## Math

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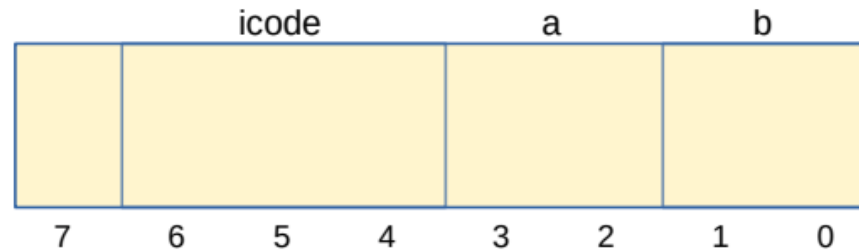
Broadly doing work

icode	b	meaning
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
6	1	$rA \&= \text{read from memory at pc} + 1$
	2	$rA += \text{read from memory at pc} + 1$

*Note: We can implement other operations using these things!*

## icodes 3 and 6

Special property of icodes 3 & 6: only one register used



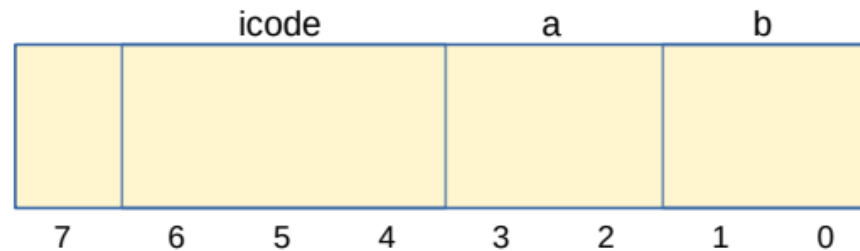
icode	b	action
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$



## icodes 3 and 6

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Special property of icodes 3 & 6: only one register used

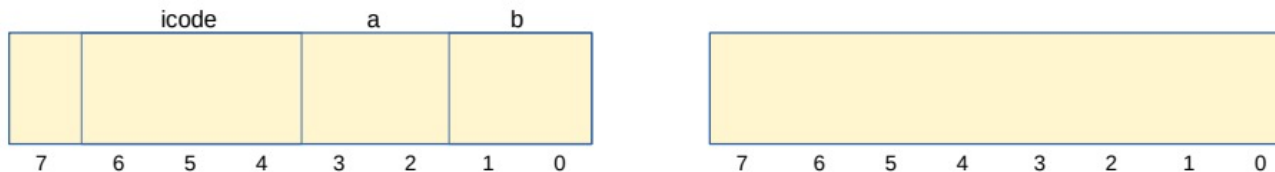


- Side effect: all bytes between 0 and 127 are valid instructions!
- As long as high-order bit is 0
- No syntax errors, any instruction given is valid

## Immediate values

icode 6 provides literals, **immediate** values

icode	b	action
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$
		For icode 6, increase pc by 2 at end of instruction



## Encoding Instructions

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Example 1:  $r1 += 19$

## Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		write $rA$ to memory at address $rB$
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$ For icode 6, increase $pc$ by 2 at end of instruction
7		Compare $rA$ as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment $pc$ as normal

## Encoding Instructions

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Example 2:  $M[0x82] += r3$

Read memory at address 0x82, add r3, write back to memory at same address

## Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		write $rA$ to memory at address $rB$
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$
		For icode 6, increase $pc$ by 2 at end of instruction
7		Compare $rA$ as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment $pc$ as normal

## Jumps

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- Moves and math are large portion of our code
- We also need **control constructs**
  - Change what we are going to do next
  - if, while, for, functions, ...
- Jumps provide mechanism to perform these control constructs
- We jump by assigning a new value to the program counter PC

## Jumps

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- For example, consider an if



## Jumps

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icode	meaning
7	Compare rA as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment pc as normal

Instruction icode 7 provides a **conditional** jump

- Real code will also provide an **unconditional** jump, but a conditional jump is sufficient

## Writing Code

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We can now write any\* program!

- When you run code, it is being turned into instructions like ours
- Modern computers use a larger pool of instructions than we have (we will get there)

\*we do have some limitations, since we can only represent 8-bit values and some operations may be tedious.

## Our code to this machine code

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How do we turn our control constructs into jump statements?

## **if/else to jump**

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## **while to jump**

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## Function Calls

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## Encoding Instructions

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Example 3: if  $r0 < 9$  jump to 0x42

## Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		write $rA$ to memory at address $rB$
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$
		For icode 6, increase $pc$ by 2 at end of instruction
7		Compare $rA$ as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment $pc$ as normal