Instruction Set Architectures, Stacks

CS 2130: Computer Systems and Organization 1 October 6, 2025

Announcements

- Homework 4 available today due Friday at 11:59pm on Gradescope
 - Note the earlier deadline!
 - You have written most of this code already
 - Lab 6 may provide a fast way to get started

Instruction Set Architecture (ISA) is an abstract model of a computer defining how the CPU is controlled by software

- Conceptually, set of instructions that are possible and how they should be encoded
- · Results in many different machines to implement same ISA
 - Example: How many machines implement our example ISA?
- Common in how we design hardware

Instruction Set Architecture (ISA) is an abstract model of a computer defining how the CPU is controlled by software

- · Provides an abstraction layer between:
 - Everything computer is really doing (hardware)
 - What programmer using the computer needs to know (software)
- Hardware and Software engineers have freedom of design, if conforming to ISA
- · Can change the machine without breaking any programs

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CSO: covering many of the times we'll need to think across this barrier

Backwards compatibility

- Include flexibility to add additional instructions later
- · Original instructions will still work
- Same program can be run on PC from 10+ years ago and new PC today

Most manufacturers choose an ISA and stick with it

· Notable Exception: Apple

What about real ISAs?

What about our ISA?

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ToyISA Instructions

icode	b	meaning
0		rA = rB
1		rA &= rB
2		rA += rB
3	0	rA = ~rA
	1	rA = !rA
	2	rA = -rA
	3	rA = pc
4		rA = read from memory at address rB
5		write rA to memory at address rB
6	0	rA = read from memory at pc + 1
	1	rA &= read from memory at pc + 1
	2	rA += read from memory at pc + 1
	3	rA = read from memory at the address stored at pc + 1
	_	For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0
		if rA <= 0 set pc = rB
		else increment pc as normal

What about our ISA?

- Enough instructions to compute what we need
- · As is, lot of things that are painful to do
 - This was on purpose! So we can see limitations of ISAs early
- · Add any number of new instructions using the reserved bit (7)
- · Missing something important: Help to put variables in memory

Storing Variables in Memory

So far... we/compiler chose location for variable Consider the following example:

```
f(x):
a = x
if (x <= 0) return 0
else return f(x-1) + a</pre>
```

Recursion

The formal study of a function that calls itself

Storing Variables in Memory

```
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a = x
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else return f(x-1) + a</pre>
```

Where do we store a?

The Stack

Stack - a last-in-first-out (LIFO) data structure

- The solution for solving this problem
- rsp Special register the stack pointer
 - Points to a special location in memory
 - Two operations most ISAs support:
 - push put a new value on the stack
 - pop return the top value off the stack

The Stack: Push and Pop

push r0

Put a value onto the "top" of the stackrsp -= 1M[rsp] = r0

pop r2

Read value from "top", save to registerr2 = M[rsp]rsp += 1

The Stack: Push and Pop

The Stack: Push and Pop

Function Calls

A short aside...

Time to take over the world!

Backdoors

Backdoor: secret way in to do new *unexpected* things

- Get around the normal barriers of behavior
- Ex: a way in to allow me to take complete control of your computer

Backdoors

Exploit - a way to use a vulnerability or backdoor that has been created

- Our exploit today: a malicious payload
 - A passcode and program
 - If it ever gets in memory, run my program regardless of what you want to do

Our backdoor will have 2 components

- · Passcode: need to recognize when we see the passcode
- Program: do something bad when I see the passcode

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- · We're talking adding a few hundred transistors

Will you notice this on your chip?

- Modern chips have billions of transistors
- We're talking adding a few hundred transistors
- Maybe with a microscope? But you'd need to know where to look!

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- · Sounds like something from the movies
- · People claim this might be happening

Have you heard about something like this before?

- Sounds like something from the movies
- People claim this might be happening
- To the best of my knowledge, no one has ever admitted to falling in this trap

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· No technical reason not to, it's easy to do!

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Can we make a system where one bad actor can't break it?

· Code reviews, double checks, verification systems, automated verification systems, ...

Why does this work?





Why?

Why does this work?

- · It's all bytes!
- Everything we store in computers are bytes
- · We store code and data in the same place: memory

It's all bytes

Memory, Code, Data... It's all bytes!

- Enumerate pick the meaning for each possible byte
- Adjacency store bigger values together (sequentially)
- Pointers a value treated as address of thing we are interested in

Enumerate

Enumerate - pick the meaning for each possible byte

What is 8-bit 0x54?

Unsigned integer Signed integer Floating point w/ 4-bit exponent ASCII Bitvector sets Our example ISA eighty-four positive eighty-four twelve capital letter T: T The set {2, 3, 5} Flip all bits of value in r1

Adjacency

Adjacency - store bigger values together (sequentially)

- An array: build bigger values out of many copies of the same type of small values
 - Store them next to each other in memory
 - Arithmetic to find any given value based on index

Adjacency

Adjacency - store bigger values together (sequentially)

- Records, structures, classes
 - Classes have fields! Store them adjacently
 - Know how to access (add offsets from base address)
 - If you tell me where object is, I can find fields

Pointers

Pointers - a value treated as address of thing we are interested in

- · A value that really points to another value
- Easy to describe, hard to use properly
- · We'll be talking about these a lot in this class!

Pointers

Pointers - a value treated as address of thing we are interested in

- Give us strange new powers (represent more complicated things), e.g.,
 - Variable-sized lists
 - Values that we don't know their type without looking
 - Dictionaries, maps

Programs Use These!

How do our programs use these?

- Enumerated icodes, numbers
- Ajacently stored instructions (PC+1)
- Pointers of where to jump/goto (addresses in memory)