



Toy Instruction Set Architecture

CS 2130: Computer Systems and Organization 1
September 26, 2025

Announcements

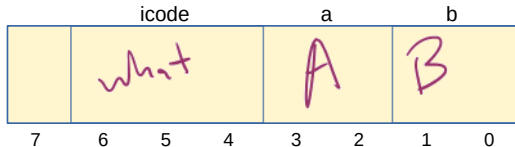
Wednesday

- Homework 3 due ~~Monday~~ at 11:59pm on Gradescope
- Quiz 4 available today, due Sunday at 11:59pm
- Midterm 1 next Friday (October 3, 2025) in class
 - Written, closed notes
 - If you have SDAC, please schedule ASAP

Encoding Instructions

Encoding of Instructions

- 3-bit icode (which operation to perform)
 - Numeric mapping from icode to operation
- Which registers to use (2 bits each)
- Reserved bit for future expansion



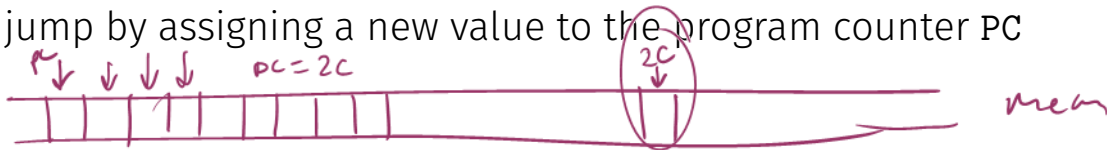
High-level Instructions

In general, 3 kinds of instructions

- **moves** - move values around without doing “work”
- **math** - broadly doing “work”
- **jumps** - jump to a new place in the code

Jumps

- Moves and math are large portion of our code
- We also need **control constructs**
 - Change what we are going to do next
 - if, while, for, functions, ...
- Jumps provide mechanism to perform these control constructs
- We jump by assigning a new value to the program counter PC



Jumps

icode	meaning
7	Compare rA as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment pc as normal

Handwritten notes in red:

- $if(\overline{rA} \leq 0) \text{ jmp } rB$
- $7 \rightarrow 0100 \rightarrow 4$

Instruction icode 7 provides a **conditional** jump

- Real code will also provide an **unconditional** jump, but a conditional jump is sufficient

Writing Code

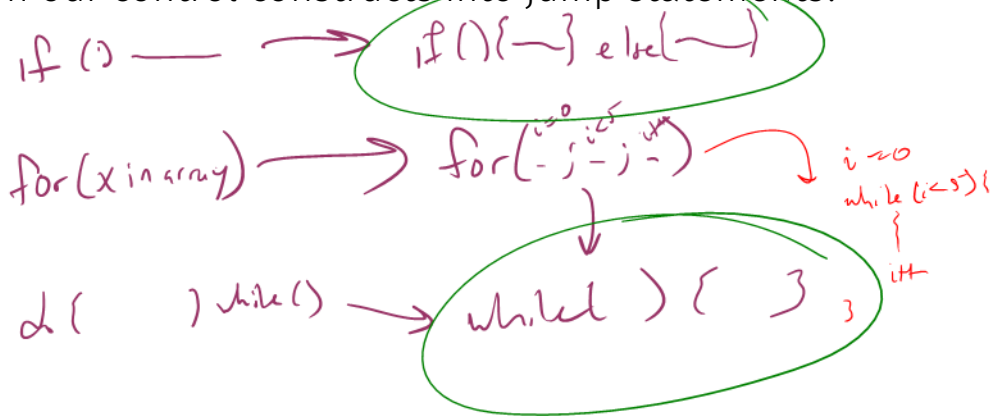
We can now write any* program!

- When you run code, it is being turned into instructions like ours
- Modern computers use a larger pool of instructions than we have (we will get there)

*we do have some limitations, since we can only represent 8-bit values and some operations may be tedious.

Our code to this machine code

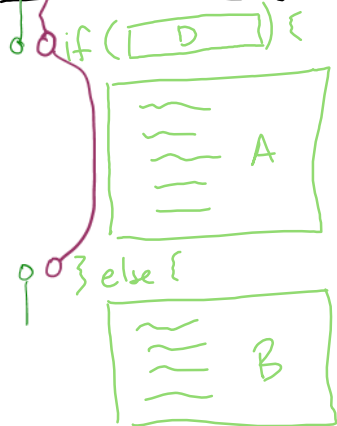
How do we turn our control constructs into jump statements?



if/else to jump

$\text{if}(rA \leq 0) \text{ jump } __$

Pseudocode using if/else

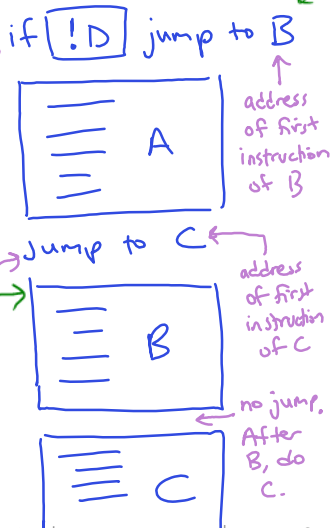


Notes!

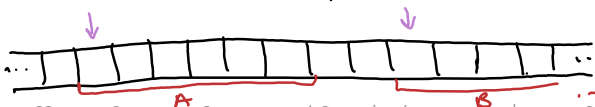
if D is true:
run code in A, then C
→ skip B

if D is false:
run code in B, then C
→ skip A

Using Jumps



Code is in memory (think array)

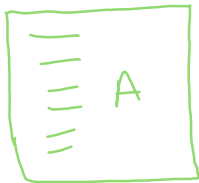


while to jump

$$3-3 = rA = PC$$

Pseudocode of while loop:

while (C) {



}



Code after loop

Notes:

→ if C is true, run code in A, then go back and check C again

if C is false, skip A, go to B

We have two options!

- jump to the check of C

- check at end of A before jumping back

option 1

address of first instruction of ...
↓

option 2

D → if (!C) jump to B



jump to D unconditionally



if (!C) jump to B



if (C) jump to A



Encoding Instructions

Example 3: if r0 < 9 jump to 0x42

```
if(x < 9) {  
    }  
    else {  
        0x42  
    }  
}
```

Instructions

-128 --- 127

-8 -3 -2 -1
F8 - - - FD FE FF 0

icode	b	meaning
0		rA = rB
1		rA &= rB
2		rA += rB
3	0	rA = ~rA
	1	rA = !rA
	2	rA = -rA
	3	rA = pc
4		rA = read from memory at address rB
5		write rA to memory at address rB
6	0	rA = read from memory at pc + 1
	1	rA &= read from memory at pc + 1
	2	rA += read from memory at pc + 1
	3	rA = read from memory at the address stored at pc + 1
		For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0 if rA <= 0 set pc = rB else increment pc as normal

$r0 < 9$
 $r0 \leq 8$
 $r0 - 8 \leq 0$

$r0 - 8 < 0$
 $r0 += -8$
 $r0 += F8$

if (r0 < 9) jump to 0x42

$r1 = 0x42$ ← 0100 / 64 42

$r0 += F8$ ← 62 F8 / 00 10

if (rA = 0) pc = r1

71 / 00 01 / A B

64 42 62 F8 71

Memory

What kinds of things do we put in memory?

- Code: binary code like instructions in our example ISA
 - Intel/AMD compatible: x86_64
 - Apple Mx and Ax, ARM: ARM
 - And others!
- Variables: we may have more variables that will fit in registers
- Data Structures: organized data, collection of data
 - Arrays, lists, heaps, stacks, queues, ...

Dealing with Variables and Memory

What if we have many variables? Compute: $x += y$

$x = 0x80$

$y = 0x81$

$z = 0x82$

$t = 0x83$

$w = 0x84$

$u = 0x85$

read from
mem

$r1 = M[0x80]$

$r2 = M[0x81]$

execute

$r1 += r2$

write to
mem

$M[0x80] = r1$

$M[0x81] = r2$

67 80 6B 81 26 60 80 54 60 81 58

$\frac{1}{01} \frac{3}{11}$
67 80

6B 81
1011

2,6
0110
60 80

54
0100
60 81
58
1000