



# Toy Instruction Set Architecture

CS 2130: Computer Systems and Organization 1  
September 24, 2025

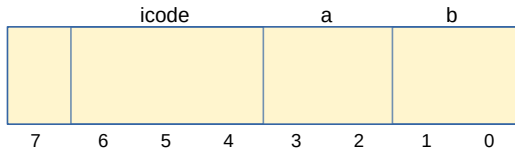
# Announcements

- Homework 3 due Monday at 11:59pm on Gradescope
- Midterm 1 next Friday (October 3, 2025) in class
  - Written, closed notes
  - If you have SDAC, please schedule ASAP
- No lab check-off on Mondays

# Encoding Instructions

## Encoding of Instructions

- 3-bit icode (which operation to perform)
  - Numeric mapping from icode to operation
- Which registers to use (2 bits each)
- Reserved bit for future expansion



# Toy ISA Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		write $rA$ to memory at address $rB$
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$
		For icode 6, increase $pc$ by 2 at end of instruction
7		Compare $rA$ as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment $pc$ as normal

# High-level Instructions

In general, 3 kinds of instructions

- **moves** - move values around without doing “work”
- **math** - broadly doing “work”
- **jumps** - jump to a new place in the code

# Moves

icode	b	action
0		$rA = rB$
3	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		$\text{write } rA \text{ to memory at address } rB$
6	0	$rA = \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$

# Math

Broadly doing work

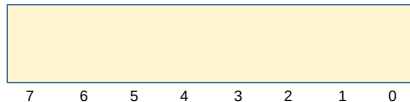
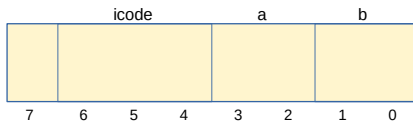
icode	b	meaning
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
6	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$

*Note: We can implement other operations using these things!*

# Immediate values

icode 6 provides literals, **immediate** values

icode	b	action
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$
		For icode 6, increase $pc$ by 2 at end of instruction





# Encoding Instructions

Example 1: `r1 += 19`

# Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		write $rA$ to memory at address $rB$
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$
		For icode 6, increase $pc$ by 2 at end of instruction
7		Compare $rA$ as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment $pc$ as normal

# Encoding Instructions

Example 2: `M[0x82] += r3`

Read memory at address 0x82, add `r3`, write back to memory at same address

# Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		write $rA$ to memory at address $rB$
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$
		For icode 6, increase $pc$ by 2 at end of instruction
7		Compare $rA$ as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment $pc$ as normal

# Writing Code: Homework Hints

1. Write pseudocode that does the desired task
- 2-3 ... deal with control flow
4. Split multi-operation lines into series of single-operation lines  
`x = y-z;` becomes `x = y; x -= z;`
5. Convert operations to those in our instruction set  
`x -= z;` becomes `w = z; w = -w; x += w;`
6. ... deal with loops
7. Assign variables to our four registers, ex: `r0=x, r1=y, r2=z, r3=w`  
`r0 = r1; r3 = r2; r3 = -r3; r0 += r3`
- 10- Write those instructions into triples, then hex

# Jumps

- Moves and math are large portion of our code
- We also need **control constructs**
  - Change what we are going to do next
  - `if`, `while`, `for`, functions, ...
- Jumps provide mechanism to perform these control constructs
- We jump by assigning a new value to the program counter PC

# Jumps

icode	meaning
7	Compare <code>rA</code> as 8-bit 2's-complement to 0 if <code>rA</code> $\leq$ 0 set <code>pc</code> = <code>rB</code> else increment <code>pc</code> as normal

Instruction icode 7 provides a **conditional** jump

- Real code will also provide an **unconditional** jump, but a conditional jump is sufficient

# Writing Code

We can now write any\* program!

- When you run code, it is being turned into instructions like ours
- Modern computers use a larger pool of instructions than we have (we will get there)

\*we do have some limitations, since we can only represent 8-bit values and some operations may be tedious.



# Our code to this machine code

How do we turn our control constructs into jump statements?

# if/else to jump

# while to jump

# Encoding Instructions

Example 3: `if r0 < 9 jump to 0x42`

# Instructions

icode	b	meaning
0		$rA = rB$
1		$rA \&= rB$
2		$rA += rB$
3	0	$rA = \sim rA$
	1	$rA = !rA$
	2	$rA = -rA$
	3	$rA = pc$
4		$rA = \text{read from memory at address } rB$
5		write $rA$ to memory at address $rB$
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA \&= \text{read from memory at } pc + 1$
	2	$rA += \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$
		For icode 6, increase $pc$ by 2 at end of instruction
7		Compare $rA$ as 8-bit 2's-complement to 0 if $rA \leq 0$ set $pc = rB$ else increment $pc$ as normal

# Function Calls